Solution of HW3

Chapter 5 Review Question R1

Although each link guarantees that an IP datagram sent over the link will be received at the other end of the link without errors, it is not guaranteed that IP datagrams will arrive at the ultimate destination in the proper order. With IP, datagrams in the same TCP connection can take different routes in the network, and therefore arrive out of order. TCP is still needed to provide the receiving end of the application the byte stream in the correct order. Also, IP can lose packets due to routing loops or equipment failures.

Chapter 5 Review Question R2

Framing: there is also framing in IP and TCP; link access; reliable delivery: there is also reliable delivery in TCP; flow control: there is also flow control in TCP; error detection: there is also error detection in IP and TCP; error correction; full duplex: TCP is also full duplex.

Chapter 5 Review Question R4

Slotted Aloha: 1, 2 and 4 (slotted ALOHA is only partially decentralized, since it requires the clocks in all nodes to be synchronized). Token ring: 1, 2, 3, 4.

Chapter 5 Review Question R6

When a node transmits a frame, the node has to wait for the frame to propagate around the entire ring before the node can release the token. Thus, if L/R is small as compared to t_{prop} , then the protocol will be inefficient.

Chapter 5 Review Question R7

 2^{48} MAC addresses; 2^{32} IPv4 addresses; 2^{128} IPv6 addresses.

Chapter 5 Problem 2

Suppose we begin with the initial two-dimensional parity matrix:

 $\begin{array}{c} 0 \ 0 \ 0 \ 0 \\ 1 \ 1 \ 1 \ 1 \\ 0 \ 1 \ 0 \ 1 \end{array}$

1010

With a bit error in row 2, column 3, the parity of row 2 and column 3 is now wrong in the matrix below:

 $\begin{array}{c} 0 \ 0 \ 0 \ 0 \\ 1 \ 1 \ 0 \ 1 \\ 0 \ 1 \ 0 \ 1 \\ 1 \ 0 \ 1 \ 0 \\ \end{array}$

Now suppose there is a bit error in row 2, column 2 and column 3. The parity of row 2 is now correct! The parity of columns 2 and 3 is wrong, but we can't detect in which rows the error occurred!

 $\begin{array}{c} 0 \ 0 \ 0 \ 0 \\ 1 \ 0 \ 1 \\ 0 \ 1 \ 0 \ 1 \\ 1 \ 0 \ 1 \\ 0 \ 1 \ 0 \\ \end{array}$

The above example shows that a double bit error can be detected (if not corrected).

Chapter 6 Review Question R2

a) Single hop, infrastructure-basedb) Single hop, infrastructure-less

c) Multi-hop, infrastructure-based

d) Multi-hop, infrastructure-less

Chapter 6 Review Question R4

a) Increasing the transmission powerb) Reducing the transmission rate

Chapter 6 Review Question R16

No. A node can remain connected to the same access point throughout its connection to the Internet (hence, not be mobile). A mobile node is the one that changes its point of attachment into the network over time. Since the user is always accessing the Internet through the same access point, she is not mobile.

Chapter 6 Review Question R17

A permanent address for a mobile node is its IP address when it is at its home network. A careof-address is the one its gets when it is visiting a foreign network. The COA is assigned by the foreign agent (which can be the edge router in the foreign network or the mobile node itself).